# Super-resolution of Positive Sources

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Joint work with Emmanuel Candès

# Diffraction limits resolution:







Ernst Abbe

# Diffraction limits resolution: $\lambda_c = \frac{\lambda_{\text{LIGHT}}}{2n\sin(\theta)}$





#### [picture from nobelprize.org]

# Looking inside the cell: conventional microscopy



microtubule

# Nobel Prize in Chemistry 2014





#### Eric Betzig

Stefan W. Hell

W.E. Moerner

Invention of single-molecule microscopy

# Looking inside the cell



#### conventional microscopy



#### single-molecule microscopy

# Single molecule microscopy (basics)



Green fluorescent protein (GFP)

Energy states [Dickson et.al. '97]



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Green fluorescent protein (GFP) Energy states [Dickson et.al. '97]

- $\blacksquare$  State A is excited to  $A^*$  and returns to A upon photon emission
- When *I* is reached from *A* there is no fluorescence until *I* spontaneously moves to *A* (blinking)
- $\blacksquare$  When I moves to N there is no fluorescence until N is activated by  $405 \mathrm{nm}$  light and GFP returns to A

# Photoactivated localization microscopy (PALM) Setup



[picture from ZEISS]

# PALM Process



#### Step 3. Algorithm needed.



#### Step 2



Step 4



Photobleach & Record Positions

# Antibodies: attach fluorescent molecules to the structure



All off

# Antibodies: attach fluorescent molecules to the structure



All off

All on

# Antibodies: attach fluorescent molecules to the structure





Frame 1



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Locate centers of "Gaussian" blobs (parametric estimation)



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Combine  $\sim$  10000 frames.



#### Locate centers of "Gaussian" blobs (parametric estimation)

Combine  $\sim$  10000 frames.

The structure is now resolved!

Imaging  $\sim 10000$  frames is  ${\rm slow}$ 

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Can we make data acquisition faster?

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Image  $\sim 2500$  frames with 4 times more molecules per frame?



parametric estimation works



4 times more active molecules  $\Rightarrow$  parametric estimation **does not** work

Imaging  $\sim 10000$  frames is  ${\rm slow}$ 

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parametric estimation works



4 times more active molecules ⇒ parametric estimation **does not** work

Need powerful super-resolution algorithm!



# Theory



# Theory

Which algorithm? Performance guarantees? Fundamental limits?

# Mathematical model (discrete 1D setup for simplicity)



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$$\mathbf{x} = [x_0 \cdots x_{N-1}]^\mathsf{T} \ge \mathbf{0}$$

$$s = Px + z$$



15/50

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# Super-resolution factor and stability

$$\mathbf{x} = [x_0 \cdots x_{N-1}]^{\mathsf{T}}$$



 $\mathbf{s} = \mathbf{P}\mathbf{x} + \mathbf{z}$ 

 $\text{SRF} \triangleq N/(2f_c)$ 

# Super-resolution factor and stability

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 $\mathbf{s} = \mathbf{P}\mathbf{x} + \mathbf{z}$ 

 $\text{SRF} \triangleq N/(2f_c)$ 

Stability: 
$$\|\mathbf{x} - \hat{\mathbf{x}}\| \stackrel{?}{\leq} \|\mathbf{z}\| \cdot (\text{amplification factor})$$

# Classical resolution criteria: separation is about $\lambda_c$



# Rayleigh-regularity: $\mathbf{x} \in \mathcal{R}(d, r)$

**x** has fewer than r spikes in every  $\lambda_c d$  interval  $[\lambda_c \triangleq 1/f_c]$
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[Prony'1795]  $\mathbf{x} \in \mathbb{C}^N$ no stability efficient













[Prony'1795]	MUSIC, ESPRIT	[Donoho'92]
$\mathbf{x} \in \mathbb{C}^N$	$\mathbf{x} \in \mathbb{C}^N$	$\mathbf{x} \in \mathbb{C}^N$
no stability	stability not understood	stability
-	-	Rayleigh-regularity
efficient	efficient	combinatorial
$\begin{array}{l} \textbf{[Donoho et al.'90]} \\ \mathbf{x} \geq 0 \\ \textbf{no stability} \\ - \\ \textbf{convex} \end{array}$	[Candès & FGranda'12] $\mathbf{x} \in \mathbb{C}^N$ stability separation convex	

Works:



[Prony'1795]	MUSIC, ESPRIT	[Donoho'92]
$\mathbf{x} \in \mathbb{C}^N$	$\mathbf{x}\in\mathbb{C}^{N}$	$\mathbf{x} \in \mathbb{C}^N$
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[Donoho et al.'90]	[Candès & FGranda'12]	
$\mathbf{x} \geq 0$	$\mathbf{x}\in\mathbb{C}^{N}$	
no stability	stability	
-	separation	
convex	convex	
Brooks		



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no stability	stability	stability
-	separation	Rayleigh-regularity
convex	convex	convex
	$\mathcal{R}(2r,r) \mathbf{v} > 0$	
	$\mathcal{R}(2l, l), \mathbf{x} \geq 0$	

#### Main results



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Converse: [V. Morgenshtern and E. Candès, 2014]

For  $\mathbf{P} = \mathbf{P}_{\text{tri}}$ , no algorithm can do better than  $c \cdot \|\mathbf{z}\|_1 \cdot \left(\frac{N}{2f_c}\right)^{2r-1}$ .

# Key ideas

- → **Duality theory:** to prove stability we need a low-frequency trigonometric polynomial that is "curvy"
  - [Dohono, et al.'92] construct trigonometric polynomial that is not "curvy"
  - [Candès and Fernandez-Granda'12] construct trigonometric polynomial that is "curvy", but construction needs separation
  - New construction: multiply "curvy" trigonometric polynomials
    - "curvy"
    - construction needs no separation

#### Dual certificate (noiseless case, z = 0)

**\blacksquare**  $\mathcal{T}$  is the support of  $\mathbf{x}$ 

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Suppose, we can construct a **low-frequency trig. polynomial**:



Then,  $\hat{\mathbf{x}} = \mathbf{x}$ .









- $\mathbf{\hat{x}} = \mathbf{x} \text{ iff there exists} \\ \mathbf{q} \perp \text{null}(\mathbf{P}) \text{ and } \mathbf{q} \in \partial \|\mathbf{x}\|_1$
- **P** is orthogonal projection onto the set of low-freq. trig. polynomials:  $\mathbf{q} \perp \text{null}(\mathbf{P}) \Leftrightarrow$  $q(t) = \sum_{k=-f_c}^{f_c} \hat{q}_k e^{-i2\pi kt}$





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$$\mathbf{q} \in \partial \|\mathbf{x}\|_1 \Leftrightarrow \\ \begin{cases} q(t_i) = \operatorname{sign}(x_i) & x_i \neq 0 \\ |q(t_i)| \le 1 & x_i = 0 \end{cases}$$





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Suppose, we can construct a **low-frequency trig. polynomial**:



■ Then,  $\|\hat{\mathbf{x}} - \mathbf{x}\|_1 \le 4 \|\mathbf{z}\|_1 / \rho$ .

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[Dohono, et al.'92]: "Classical" q(t)

$$q(t) = \prod_{t_0 \in \mathcal{T}} \frac{1}{2} \left[ \cos(2\pi(t+1/2-t_0)) + 1 \right].$$

No separation required

Low curvature!

$$q(t-t_0) \approx (t-t_0)^2 \Rightarrow \|\mathbf{x} - \hat{\mathbf{x}}\|_1 \le \|\mathbf{z}\|_1 \cdot N^2$$



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#### [Candès, Fernandez-Granda'12]: "Curvy" q(t)

$$q(t) = \sum_{t_j \in \mathcal{T}} a_j K(t-t_j) + \text{corrections},$$
 
$$K(t) \dots \text{low-frequency and "curvy"}$$

Separation between zeros required:  $\mathcal{T} \in \mathcal{R}(2,1)$ 

High curvature!

$$q(t-t_i) \approx f_c^2 (t-t_i)^2 \Rightarrow \|\mathbf{x} - \hat{\mathbf{x}}\|_1 \le c \cdot \|\mathbf{z}\|_1 \cdot \left(\frac{N}{2f_c}\right)^2$$



#### Comparison of Trigonometric Polynomials



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#### New construction: curvature without separation

Partition support:  $\mathcal{T} = \mathcal{T}_1 \cup \mathcal{T}_2$ , r = 2Regularity:  $\mathcal{T} \in \mathcal{R}(2 \cdot 2, 2) \Rightarrow \mathcal{T}_i \in \mathcal{R}(4, 1)$ 



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 $q(t; f_c) = q_1(t; f_c/2) \times q_2(t; f_c/2)$ 

#### High curvature!

$$q(t-t_i) \approx \frac{f_c^{2r}}{r^{2r}} (t-t_i)^{2r} \Rightarrow \|\mathbf{x} - \hat{\mathbf{x}}\|_1 \le c \cdot \|\mathbf{z}\|_1 \cdot \left(\frac{N}{2f_c}\right)^{2r}$$

Remember: q(t) must be frequency-limited to  $f_c!$ 

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[Donoho, et.al.]:

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This work:

$$q(t) = \prod_{k=1}^{r} \sum_{t_{jk} \in \mathcal{T}_k} \underbrace{a_{jk} K(t - t_{jk})}_{\text{frequency } f_c/r}$$

#### Complex vs. positive signals

Why do we need  $x \ge 0$ ?

 $\mathbf{x} \ge \mathbf{0}$   $\mathbf{x} \in \mathbb{C}^N$ 

Interpolate **zero** on supp. of  $\mathbf{x}$ 

Interpolate  $\mathrm{sign}(\mathbf{x})$  on supp. of  $\mathbf{x}$ 





Does not exist! (Bernstein Th.)
# Continuous setup

### $f_c \text{ fixed}, N \to \infty \Rightarrow \text{SRF}_{\text{OLD}} \to \infty$



Is the problem hopeless?

### $f_c \text{ fixed, } N \to \infty \Rightarrow \text{SRF}_{\text{OLD}} \to \infty$



Is the problem hopeless?

No: we need to be less ambitions!

#### $f_c \text{ fixed, } N \to \infty \Rightarrow \text{SRF}_{\text{OLD}} \to \infty$



#### $f_c \text{ fixed, } N \to \infty \Rightarrow \text{SRF}_{\text{OLD}} \to \infty$



 $\mathrm{SRF}_{\mathrm{NEW}} = \lambda_c / \lambda_{\mathrm{hi}}$ 

#### Need new tools

Theorem: [V. Morgenshtern and E. Candès, 2014] Assume  $x(t) \ge 0$ ,  $x(t) \in \mathcal{R}(2r, r)$ . Then,

$$\|f_{\mathrm{hi}} \star (x - \hat{x})\|_{1} \le c \cdot \left(\frac{\lambda_{c}}{\lambda_{\mathrm{hi}}}\right)^{2r} \cdot \|z(t)\|_{1}.$$

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Can do: all zeros

**Need:** arbitrary pattern  $\{0, +\rho\}$ 



### 2D Super-resolution



Theorem: [V. Morgenshtern and E. Candès, 2014]

Take  $\mathbf{P} = \mathbf{P}_{tri,2D}$  or  $\mathbf{P} = \mathbf{P}_{flat,2D}$ . Assume  $\mathbf{x} \ge 0$ ,  $\mathbf{x} \in \mathcal{R}(2.38r, r)$ . Then,

$$\|\hat{\mathbf{x}} - \mathbf{x}\|_1 \le c \cdot \left(\frac{N}{2f_c}\right)^{2r} \delta.$$

**New:** number of spikes is linear in the number of observations

# Improving microscopes

Collaboration with Moerner Lab, C.A. Sing-Long, E. Candès

# Reconstruction of 3D signals from 2D data





#### 2D double-helix data

# Reconstruction of 3D signals from 2D data





2D double-helix data

$$\begin{array}{ll} \mbox{minimize} & \frac{1}{2} \| \mathbf{s} - \mathbf{P} \hat{\mathbf{x}} \|_2^2 + \lambda \sigma \| \operatorname{diag}(\mathbf{w}) \hat{\mathbf{x}} \|_1 \\ \mbox{subject to} & \hat{\mathbf{x}} \geq 0 \end{array}$$

### Preliminary result: 4 times faster than state-of-the-art



10000 CVX problems solved TFOCS first order solver millions of variables

#### Flexible framework: smooth background separation



minimize subject to

$$\begin{aligned} &\frac{1}{2} \|\mathbf{s} - \mathbf{P}(\hat{\mathbf{x}} + \mathbf{b})\|_2^2 + \lambda \sigma \|\hat{\mathbf{x}}\|_1 \\ &\hat{\mathbf{x}} \ge 0 \\ &\mathbf{b} \text{ low freq. trig. polynomial (background)} \end{aligned}$$

# Comparison of super-resolution algorithms

#### Work in progress:

Need to carefully compare super-resolution algorithms in practice

- Naive matched-filters
- Algebraic methods: MUSIC, ESPRIT, ...
- Convex-optimization algorithms with different regularizers
- Realistic physical model
  - Noise: quantum noise and out-of-focus background
  - Point-spread function  $(\mathbf{P})$  uncertainty and variation
  - Rotation of single molecules

····

- Test images (phantoms):
  - Different densities of sources
  - Different spacial distributions
  - Correct answer should always be known

Create a database of test cases for quick algorithm assessment

# Convex optimization is a near-optimal method for super-resolution of positive sources

- Flexibility and good practical performance
- Non-asymptotic precise stability bounds
- Rayleigh-regularity is fundamental: separation between spikes is only one part of the picture

### Lots of questions remain

- What is the best regularizer in the presence of stochastic noise?
- Fast parallel solver exploiting the structure of the problem
- Theory for Double-Helix reconstruction: 3D signal from 2D observations
- Tractable near-optimal algorithm for complex-valued signals?

# Lots of questions remain

- What is the best regularizer in the presence of stochastic noise?
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- Theory for Double-Helix reconstruction: 3D signal from 2D observations
- Tractable near-optimal algorithm for complex-valued signals?
- Sparse regression where the design matrix has highly correlated columns:
  - $\mathbf{A} = [\mathbf{a}_1, \dots, \mathbf{a}_N]$
  - Shift-invariance:  $\langle \mathbf{a}_k, \mathbf{a}_l \rangle = \langle \mathbf{a}_{k+r}, \mathbf{a}_{l+r} \rangle$
  - $\langle \mathbf{a}_k, \mathbf{a}_{k+r} \rangle$  is large for small r
  - $\langle \mathbf{a}_k, \mathbf{a}_{k+r} 
    angle$  decays quickly with r
  - Minimum separation likely needed in general
  - If all elements in A are nonnegative and the signal is nonnegative, regularity might be enough

- Swiss National Science Foundation Fellowship
- Simons Foundation
- Theory: collaboration with E. Candès
- Motivation and applications: collaboration with W.E. Moerner, C.A. Sing-Long, M.D. Lew, A. Backer, S.J. Sahl
- Helpful discussions and related work: C. Fernandez-Granda, M. Soltanolkotabi, R. Heckel

# Thank you

# Backup slides

• Set: 
$$\mathbf{h} = \hat{\mathbf{x}} - \mathbf{x}, \quad \mathcal{T} = \{l/N : h_l < 0\}$$

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Dual vector  $q_l = q(l/N)$  satisfies:
 $\mathbf{P}_{\operatorname{flat}}\mathbf{q} = \mathbf{q}$ ,  $\|\mathbf{q}\|_{\infty} = 1$ , and  $\begin{cases} q_l = 0, \quad l/N \in \mathcal{T} \\ q_l > \rho, \quad \operatorname{otherwise.} \end{cases}$ 

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On the one hand:

$$\begin{aligned} |\langle \mathbf{q} - \rho/2, \mathbf{h} \rangle| &= |\langle \mathbf{P}(\mathbf{q} - \rho/2), \mathbf{h} \rangle| = |\langle \mathbf{q} - \rho/2, \mathbf{P}\mathbf{h} \rangle| \\ &\leq \|\mathbf{q} - \rho/2\|_{\infty} \|\mathbf{P}\mathbf{h}\|_{1} \leq \|\mathbf{P}\mathbf{x} - \mathbf{s} + \mathbf{s} - \mathbf{P}\hat{\mathbf{x}}\|_{1} \\ &\leq \|\mathbf{P}\mathbf{x} - \mathbf{s}\|_{1} + \|\mathbf{s} - \mathbf{P}\hat{\mathbf{x}}\|_{1} \\ &\leq 2\|\mathbf{P}\mathbf{x} - \mathbf{s}\|_{1} \leq 2\|\mathbf{z}\|_{1}. \end{aligned}$$

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$$|\langle \mathbf{q} - \rho/2, \mathbf{h} \rangle| = \left| \sum_{l=0}^{N-1} (q_l - \rho/2) h_l \right| = \sum_{l=0}^{N-1} (q_l - \rho/2) h_l \ge \rho \|\mathbf{h}\|_1 / 2.$$

Set: 
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• Combining:  $\|\mathbf{h}\|_1 \leq 4\|\mathbf{z}\|_1/\rho$ .

# Connection to Bernstein theorem

**Consider:** 
$$q(t) = \sum_{k=-f_c}^{f_c} \hat{q}_k e^{-i2\pi kt}$$
 with  $||q||_{\infty} \le 1$   
**Then:**  $||q'||_{\infty} \le 2f_c$ 

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Since

$$q(t_i) = 0$$
$$q'(t_i) = 0$$
$$\|q\|_{\infty} \le 1$$

We conclude:

$$\begin{split} \|q'\|_{\infty} &\leq 2f_c \Rightarrow \|q''\|_{\infty} \leq (2f_c)^2 \\ &\Rightarrow q(t-t_i) \leq (2f_c)^2 (t-t_i)^2 \\ &\Rightarrow q(t_i+1/N) \leq \frac{(2f_c)^2}{N^2} = \frac{1}{\mathrm{SRF}^2} \end{split}$$

#### New tools

Control behavior on separated set
 Multiply

$$q(t) = q_1(t) \times q_2(t)$$
  

$$0 = q'(t_3) = q'_1(t_3)q_2(t_3) + q_1(t_3)q'_2(t_3)$$



#### New tools

Control behavior on separated set
 Multiply

$$q(t) = q_1(t) \times q_2(t)$$
  

$$0 = q'(t_3) = q'_1(t_3)q_2(t_3) + q_1(t_3)q'_2(t_3)$$



3 Sum

$$q(t) = \sum_{r} \prod_{k=1}^{r} \sum_{t_{jk} \in \mathcal{T}_k} \underbrace{a_{jk} K(t - t_{jk})}_{\text{frequency } f_c/r}$$