Capacity Pre-log of Noncoherent SIMO Channels via Hironaka's Theorem

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Joint work with E. Riegler, W. Yang, G. Durisi, S. Lin, B. Sturmfels, and H. Bőlcskei











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Noncoherent setting: h_t is not known at RX

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Channel gains in whitened form:

$$\begin{bmatrix} h_1 \\ h_2 \\ \vdots \\ h_T \end{bmatrix} = \mathbf{P} \begin{bmatrix} s_1 \\ \vdots \\ s_Q \end{bmatrix}, \qquad s_q \sim \mathcal{CN}(0, 1), \text{ iid across } q$$

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Fundamental problem:

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Mutual information: $I({x_t}; {y_t}) = h({x_t}) - h({x_t} | {y_t})$ Differential Entropy: $h({y_t})$

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Pre-Log

$$\chi = \lim_{\text{SNR} \to \infty} \frac{C(\text{SNR})}{\log(\text{SNR})}$$

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Number of unknown channel parameters per block is $Q = \operatorname{rank} \mathbf{P}$: $[h_1 \ h_2 \ \dots h_T]^{\mathsf{T}} = \mathbf{P}[s_1 \ s_2 \ \dots s_Q]^{\mathsf{T}}$

Eliminating the Noncoherent Penalty by Adding Receive Antennas (Only!)



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■ Is the **pre-log** then given by (1 - MQ/T)?



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We can actually do better!

Main Result [VM et. al., 2012]

Under technical conditions on \mathbf{P} ,

$$\chi_{\text{SIMO}} = \min[1 - 1/T, R(1 - Q/T)]$$

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Multiple antennas at the receiver can recover degrees of freedom

Implications

 $M = 2, T \ge 2Q - 1$ $\chi_{\text{SIMO}} = 1 - 1/T \stackrel{!}{>} 1 - \underbrace{Q/T}_{\chi_{\text{SISO}}}$

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The "channel identification penalty" vanishes if a second receive antenna is used

Linear Algebra: Guessing Pre-Log Noiseless I/O relation:

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Rule of thumb: $\chi = \frac{\text{number of RVs } x_t \text{ that can be identified uniquely from } \{\hat{y}_{mt}\}}{T}$

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$$\begin{bmatrix} h_1 \\ h_2 \\ h_3 \end{bmatrix} = \underbrace{\begin{bmatrix} 1 & 0 \\ 0 & 1 \\ 1 & 1 \end{bmatrix}}_{\mathbf{P}} \begin{bmatrix} s_1 \\ s_2 \end{bmatrix}$$

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Quadratic equations!

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- $\Rightarrow \chi_{SISO} = (3-2)/3 = 1/3 = 1 Q/T$ [Liang & Veeravalli, 04]

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$$\Rightarrow \chi_{\text{SIMO}} = (3-1)/3 = 2/3 = 1 - 1/T$$

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$$\blacksquare \Rightarrow \chi_{\text{SIMO}} = (3-1)/3 = 2/3 = 1 - 1/T > 1/3 = \chi_{\text{SISO}}$$

Equations for General P
$$(T = 3, Q = 2, M = 2)$$



Solution is unique iff ${\bf B}$ is full-rank

Information Theory

Vector Notations
$$(T = 3, Q = 2, M = 2)$$

$$\underbrace{\begin{bmatrix} y_{11} \\ y_{12} \\ y_{13} \\ y_{21} \\ y_{22} \\ y_{23} \end{bmatrix}}_{\mathbf{y}} = \sqrt{\mathrm{SNR}} \underbrace{\begin{bmatrix} s_{11}x_1 \\ s_{12}x_2 \\ (s_{11} + s_{12})x_3 \\ s_{21}x_1 \\ s_{22}x_2 \\ (s_{21} + s_{22})x_3 \end{bmatrix}}_{\hat{\mathbf{y}}} + \underbrace{\begin{bmatrix} w_{11} \\ w_{12} \\ w_{13} \\ w_{21} \\ w_{22} \\ w_{23} \end{bmatrix}}_{\mathbf{y}}$$

 $\mathbf{P} = \begin{bmatrix} 1 & 0\\ 0 & 1\\ 1 & 1 \end{bmatrix}$

$$\mathbf{s} = [s_{11} \ s_{12} \ s_{21} \ s_{22}]^{\mathsf{T}}$$

 $\mathbf{x} = [x_1 \ x_2 \ x_3]^{\mathsf{T}}$

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- Transmit $x_i \sim \mathcal{CN}(0,1)$, i.i.d.
- $\blacksquare I(\mathbf{x}; \mathbf{y}) = h(\mathbf{y}) h(\mathbf{y} \,|\, \mathbf{x})$
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$$= h\left(\sqrt{\mathrm{SNR}}\hat{\mathbf{y}}\right)$$

$$= 6\log(\mathrm{SNR}) + h(\hat{\mathbf{y}})$$

finite?

Transmit
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- $I(\mathbf{x}; \mathbf{y}) = h(\mathbf{y}) h(\mathbf{y} \mid \mathbf{x}) \ge \frac{6}{6} \log(\text{SNR}) \frac{4}{6} \log(\text{SNR}) + c$
- **y** is Gaussian conditioned on $\mathbf{x} \Rightarrow h(\mathbf{y} \,|\, \mathbf{x}) \approx 4 \log(\text{snr})$

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$$\chi_{\text{SIMO}} \ge (\mathbf{6} - \mathbf{4})/3 = 2/3 > 1/3 = \chi_{\text{SISO}}$$

Is $h(\hat{\mathbf{y}})$ Finite? Change of Variables

$$\hat{\mathbf{y}} = \begin{bmatrix} s_{11}x_1 \\ s_{12}x_2 \\ (s_{11} + s_{12})x_3 \\ s_{21}x_1 \\ s_{22}x_2 \\ (s_{21} + s_{22})x_3 \end{bmatrix}$$
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■ $h(\hat{\mathbf{y}}) \ge h(\hat{\mathbf{y}} | x_1)$ (pilot-symbol in the noiseless case)

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Change of Variables Lemma:

$$h(\hat{\mathbf{y}} \mid x_1) = h(\mathbf{s}, x_2, x_3 \mid x_1) + \mathbb{E}_{\mathbf{s}, \mathbf{x}} \log \left| \det \frac{\partial \hat{\mathbf{y}}}{\partial(\mathbf{s}, x_2, x_3)} \right|$$

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Change of Variables Lemma:

$$h(\hat{\mathbf{y}} \mid x_1) = \underbrace{h(\mathbf{s}, x_2, x_3 \mid x_1)}_{\text{finite!}} + 2\underbrace{\mathbb{E}_{\mathbf{s}, \mathbf{x}} \log \left| \det \frac{\partial \hat{\mathbf{y}}}{\partial(\mathbf{s}, x_2, x_3)} \right|}_{\text{finite?}}$$

$$\mathbb{E}_{\mathbf{s},\mathbf{x}} \log \left| \det \frac{\partial \hat{\mathbf{y}}}{\partial (\mathbf{s}, x_2, x_3)} \right| > -\infty?$$

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Factorize:
$$\frac{\partial \hat{\mathbf{y}}}{\partial (\mathbf{s}, x_2, x_3)} = J_1(\mathbf{x}) J_2(\mathbf{s}) J_3(\mathbf{x})$$

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 $J_1(\mathbf{x})$ and $J_3(\mathbf{x})$ are diagonal matrices

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$$J_1(\mathbf{x})$$
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• det $J_2(\mathbf{s})$ is a homogeneous polynomial:

$$\det J_2(\lambda \mathbf{s}) = \lambda^D \det J_2(\mathbf{s}), \ \forall \lambda \in \mathbb{C}$$

Resolution of Singularities (Cont'd)

Polar coordinates: $\mathbf{s} \rightarrow (r, \theta)$

$$\begin{aligned} \left| \int_{\mathbb{C}^{RQ}} \exp(-\|\mathbf{s}\|^2) \log |\det J_2(\mathbf{s})| \, d\mathbf{s} \right| \\ &\leq \left| \int_{\mathbb{C}^{RQ}} \exp(-\|s\|^2) \log \left| \det J_2(\mathbf{s}/\|s\|^2) \right| \, d\mathbf{s} \right| + O(1) \\ &\leq \int_0^\infty \exp(-r^2) r^{2D-1} dr \times \int_\Delta |\log |f(\theta)|| \, d\theta + O(1) \end{aligned}$$

where

•
$$\Delta = [0,\pi]^{2D-2} \times [0,2\pi]$$
 is a compact set

• f is a real analytic function

Resolution of Singularities (Cont'd)

Hironaka's Theorem implies: If $f \not\equiv 0$ is a real analytic function, then $\int_{\Delta} |\log |f(\theta)|| \, d\theta < \infty.$



The Technical Condition on $\mathbf{P}:$ not Just Rank

	p_{11}	p_{12}	0	0	0	0]	
$f \not\equiv 0$ iff	p_{21}	p_{22}	0	0	p_{21}	0	is full-rank
	p_{31}	p_{32}	0	0	0	p_{31}	
	0	0	p_{11}	p_{12}	0	0	
	0	0	p_{21}	p_{22}	p_{22}	0	
	0	0	p_{31}	p_{32}	0	p_{32}	

The Technical Condition on $\mathbf{P}:$ not Just Rank

$$f \not\equiv 0 \text{ iff } \begin{bmatrix} p_{11} & p_{12} & 0 & 0 & 0 & 0 \\ p_{21} & p_{22} & 0 & 0 & p_{21} & 0 \\ p_{31} & p_{32} & 0 & 0 & 0 & p_{31} \\ 0 & 0 & p_{11} & p_{12} & 0 & 0 \\ 0 & 0 & p_{21} & p_{22} & p_{22} & 0 \\ 0 & 0 & p_{31} & p_{32} & 0 & p_{32} \end{bmatrix} \text{ is full-rank}$$

For T = 3, Q = 2, M = 2 equivalent to:

Every two rows of \mathbf{P} are linearly independent

The Technical Condition on $\mathbf{P}:$ not Just Rank

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For T = 3, Q = 2, M = 2 equivalent to:

Every two rows of \mathbf{P} are linearly independent

For example,
$$\mathbf{P} = \begin{bmatrix} 1 & 0 \\ 0 & 1 \\ 1 & 1 \end{bmatrix}$$
 satisfies this condition



Open problems

- MIMO
- Stationary Channel Model

Thank you