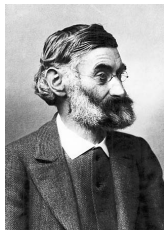
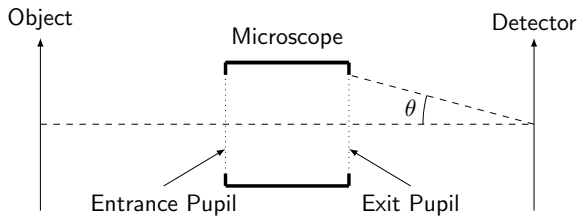


Lecture 12-13: Super-resolution of Positive Sources

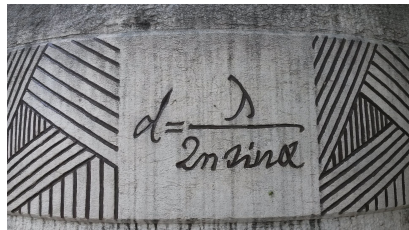
V. Morgenshtern

Mathematical Methods in Machine Learning and Signal Processing
SS 2019

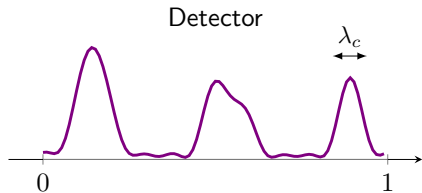
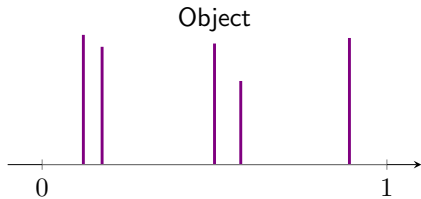
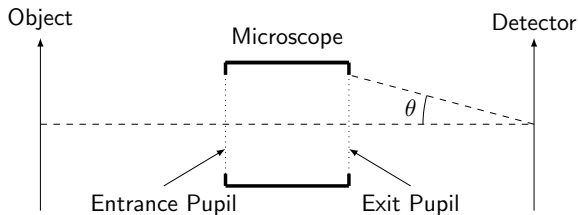
Diffraction limits resolution:

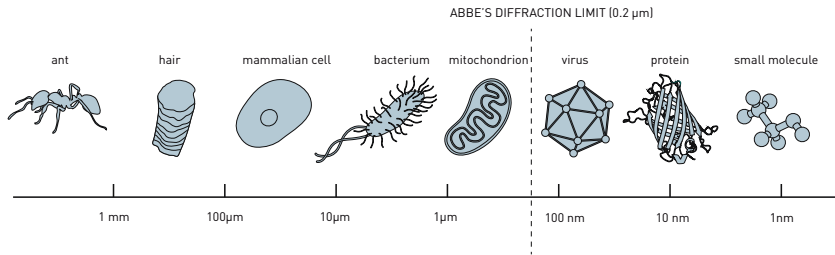


Ernst Abbe



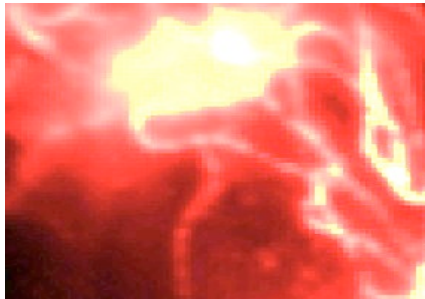
Diffraction limits resolution: $\lambda_c = \frac{\lambda_{\text{LIGHT}}}{2n \sin(\theta)}$





[picture from nobelprize.org]

Looking inside the cell: conventional microscopy



microtubule

Nobel Prize in Chemistry 2014



Eric Betzig



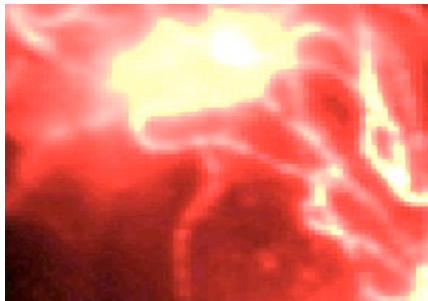
Stefan W. Hell



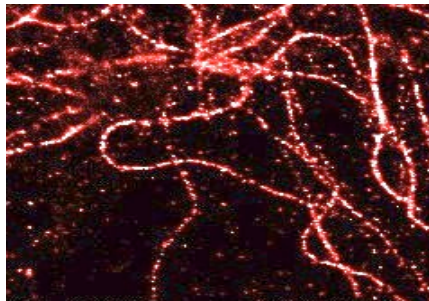
W.E. Moerner

Invention of single-molecule microscopy

Looking inside the cell



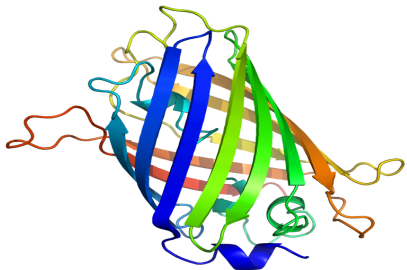
conventional microscopy



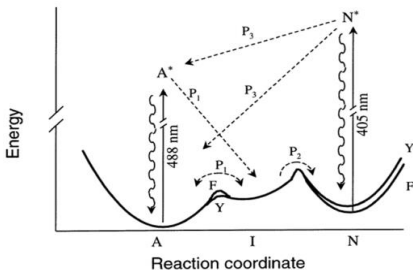
single-molecule microscopy

Single molecule microscopy (basics)

Controlled photoactivation

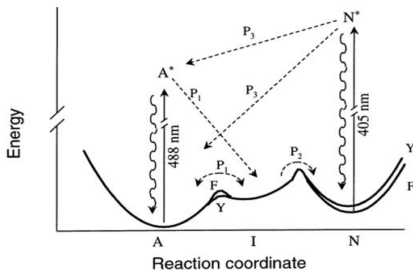
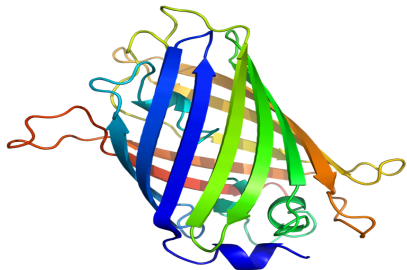


Green fluorescent protein (GFP)



Energy states [Dickson et.al. '97]

Controlled photoactivation

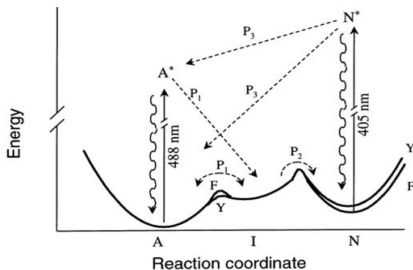
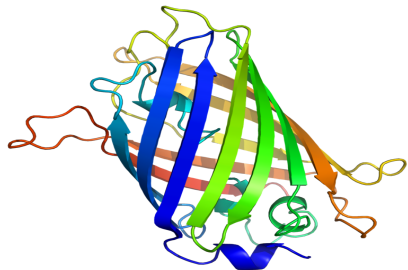


Green fluorescent protein (GFP)

Energy states [Dickson et.al. '97]

- State A is excited to A^* and returns to A upon photon emission

Controlled photoactivation

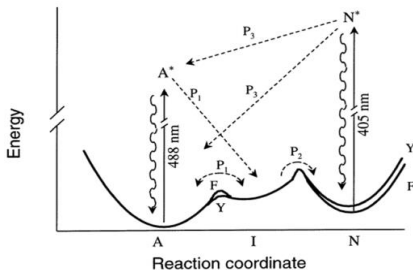
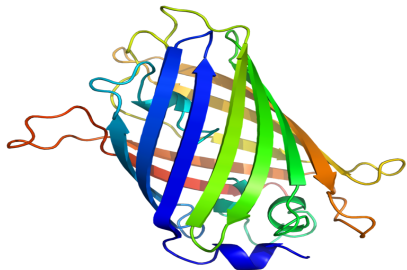


Green fluorescent protein (GFP)

Energy states [Dickson et.al. '97]

- State A is excited to A^* and returns to A upon photon emission
- When I is reached from A there is no fluorescence until I spontaneously moves to A (blinking)

Controlled photoactivation

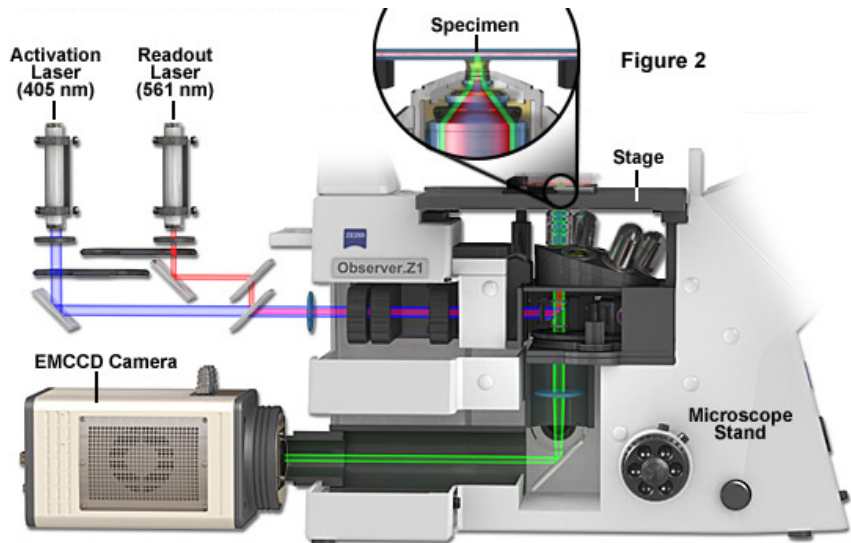


Green fluorescent protein (GFP)

Energy states [Dickson et.al. '97]

- State A is excited to A^* and returns to A upon photon emission
- When I is reached from A there is no fluorescence until I spontaneously moves to A (blinking)
- When I moves to N there is no fluorescence until N is activated by 405nm light and GFP returns to A

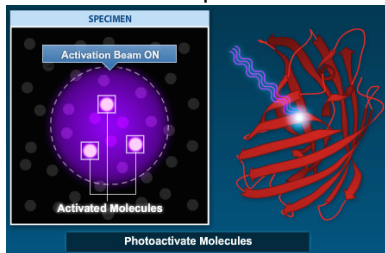
Photoactivated localization microscopy (PALM) Setup



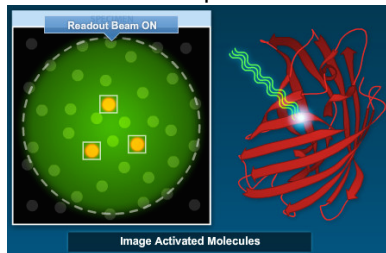
[picture from ZEISS]

PALM Process

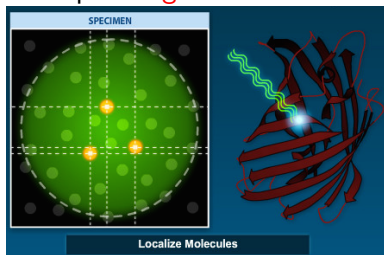
Step 1



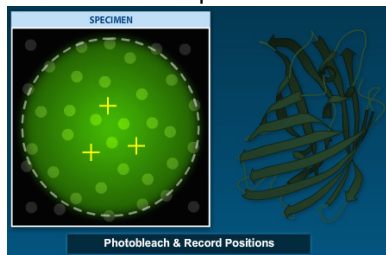
Step 2



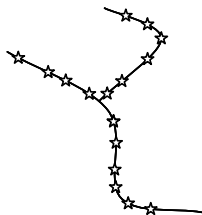
Step 3. Algorithm needed.



Step 4

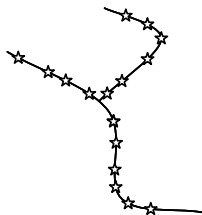


Antibodies: attach fluorescent molecules to the structure

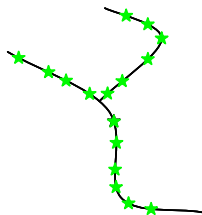


All off

Antibodies: attach fluorescent molecules to the structure

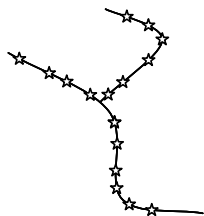


All off

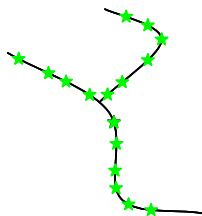


All on

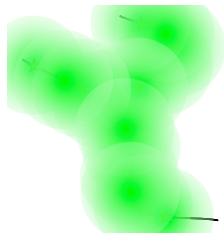
Antibodies: attach fluorescent molecules to the structure



All off



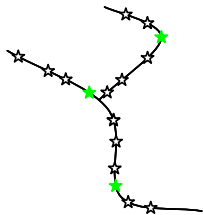
All on



Detector

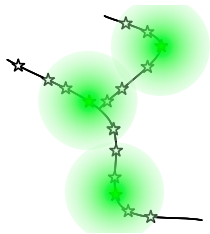
Cannot resolve the structure!

“Blinking” molecules: sparsity



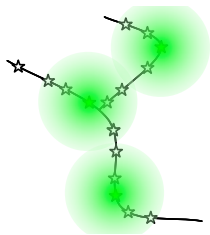
Frame 1

“Blinking” molecules: sparsity



Frame 1

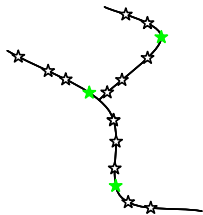
“Blinking” molecules: sparsity



Frame 1

Locate centers of “Gaussian” blobs (parametric estimation)

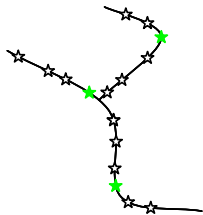
“Blinking” molecules: sparsity



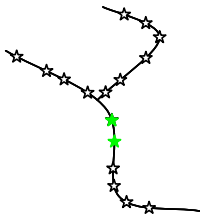
Frame 1

Locate centers of “Gaussian” blobs (parametric estimation)

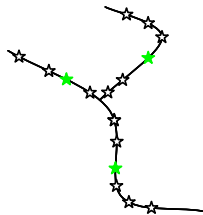
“Blinking” molecules: sparsity



Frame 1



Frame 2

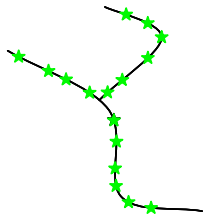


Frame 3

Locate centers of “Gaussian” blobs (parametric estimation)

Combine ~ 10000 frames.

“Blinking” molecules: sparsity



Locate centers of “Gaussian” blobs (parametric estimation)

Combine ~ 10000 frames.

The structure is now resolved!

Next Frontier: image dynamical processes

Imaging ~ 10000 frames is **slow**

Next Frontier: image dynamical processes

Imaging ~ 10000 frames is **slow**

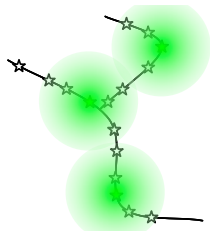
Can we make data acquisition **faster**?

Next Frontier: image dynamical processes

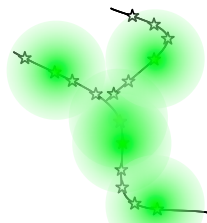
Imaging ~ 10000 frames is **slow**

Can we make data acquisition **faster**?

Image ~ 2500 frames with 4 times more molecules per frame?



parametric estimation works



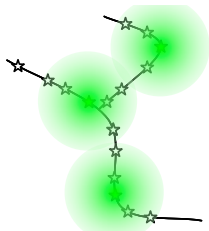
4 times more active molecules
 \Rightarrow parametric estimation
does not work

Next Frontier: image dynamical processes

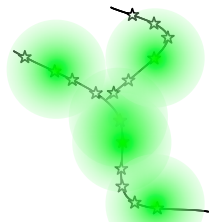
Imaging ~ 10000 frames is **slow**

Can we make data acquisition **faster**?

Image ~ 2500 frames with 4 times more molecules per frame?

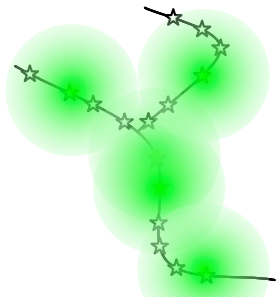


parametric estimation works

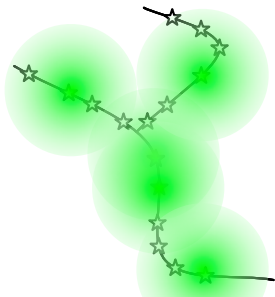


4 times more active molecules
 \Rightarrow parametric estimation
does not work

Need powerful super-resolution algorithm!



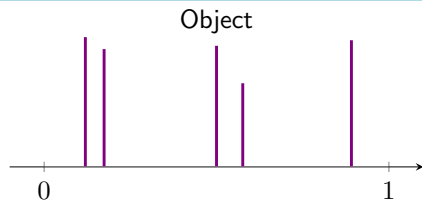
Theory



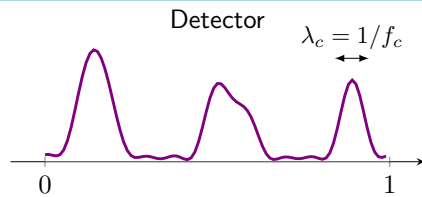
Theory

Which algorithm?
Performance guarantees?
Fundamental limits?

Mathematical model (discrete 1D setup for simplicity)



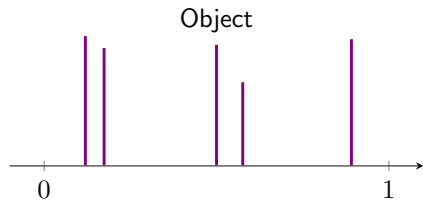
$$x(t) = \sum_i x_i \delta(t - t_i), \quad x_i \geq 0$$



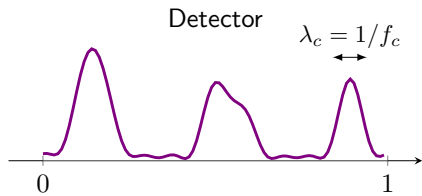
$$s(t) = \int f_{\text{low}}(t - t') x(t') dt'$$

$$f_{\text{low}}(t) = \frac{1}{2f_c} \left(\frac{\sin(2\pi f_c t)}{\pi t} \right)^2$$

Mathematical model (discrete 1D setup for simplicity)



$$x(t) = \sum_i x_i \delta(t - t_i), \quad x_i \geq 0$$

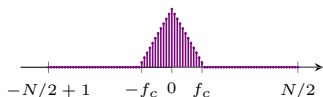


$$s(t) = \int f_{\text{low}}(t - t') x(t') dt'$$

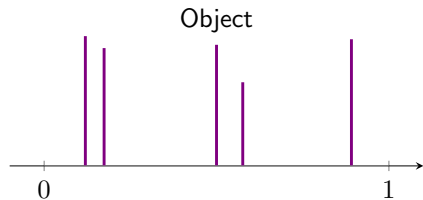
$$\mathbf{x} = [x_0 \cdots x_{N-1}]^T \geq \mathbf{0}$$

$$\mathbf{s} = \mathbf{P}\mathbf{x} + \mathbf{z}$$

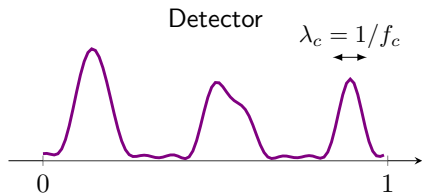
$\mathbf{P} = \mathbf{P}_{\text{tri}}$ is **circulant**
Triangular spectrum



Mathematical model (discrete 1D setup for simplicity)



$$x(t) = \sum_i x_i \delta(t - t_i), \quad x_i \geq 0$$

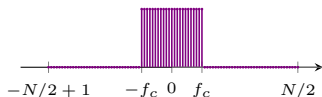


$$s(t) = \int f_{\text{low}}(t - t') x(t') dt'$$

$$\mathbf{x} = [x_0 \cdots x_{N-1}]^T \geq \mathbf{0}$$

$$\mathbf{s} = \mathbf{P}\mathbf{x} + \mathbf{z}$$

$\mathbf{P} = \mathbf{P}_{\text{flat}}$ is **circulant**
Flat spectrum



Mathematical model (discrete 1D setup for simplicity)

$$\mathbf{P} = \mathbf{F}^H \hat{\mathbf{P}} \mathbf{F}$$

DFT:

$$[\mathbf{F}]_{k,l} = \frac{1}{\sqrt{N}} e^{-i2\pi kl/N}, \quad -N/2 + 1 \leq k \leq N/2, \quad 0 \leq l \leq N - 1$$

Spectrum:

$$\hat{\mathbf{P}} = \text{diag}([\hat{p}_{-N/2+1} \cdots \hat{p}_{N/2}]^T)$$

■ **Flat:**

$$\hat{p}_k = \begin{cases} 1, & k = f_c, \dots, f_c, \\ 0, & \text{otherwise} \end{cases}$$

■ **Triangular:**

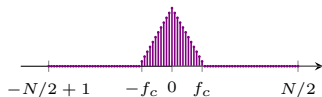
$$\hat{p}_k = \begin{cases} 1 - \frac{|k|}{f_c+1}, & k = -f_c, \dots, f_c \\ 0, & \text{otherwise} \end{cases}$$

Width of the convolution kernel: $\lambda_c \triangleq 1/f_c$

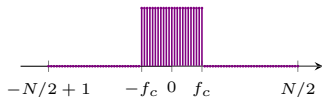
Super-resolution factor and stability

$$\mathbf{x} = [x_0 \cdots x_{N-1}]^T$$

Triangular spectrum



Flat spectrum



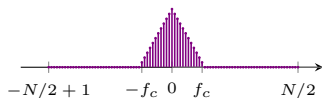
$$\mathbf{s} = \mathbf{P}\mathbf{x} + \mathbf{z}$$

$$\text{SRF} \triangleq N/(2f_c)$$

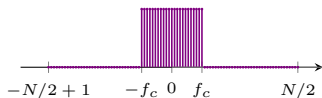
Super-resolution factor and stability

$$\mathbf{x} = [x_0 \cdots x_{N-1}]^T$$

Triangular spectrum



Flat spectrum

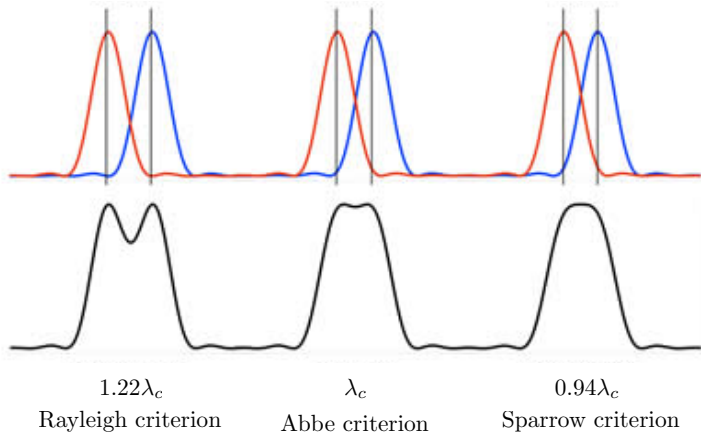


$$\mathbf{s} = \mathbf{P}\mathbf{x} + \mathbf{z}$$

$$\text{SRF} \triangleq N/(2f_c)$$

$$\text{Stability: } \|\mathbf{x} - \hat{\mathbf{x}}\| \stackrel{?}{\leq} \|\mathbf{z}\| \cdot (\text{amplification factor})$$

Classical resolution criteria: separation is about λ_c

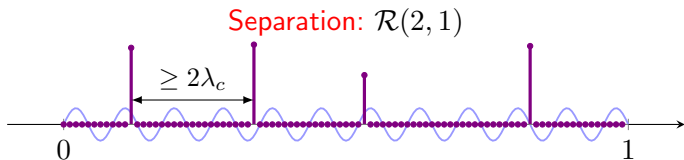


Rayleigh-regularity: $\mathbf{x} \in \mathcal{R}(d, r)$

\mathbf{x} has fewer than r spikes in every $\lambda_c d$ interval [$\lambda_c \triangleq 1/f_c$]

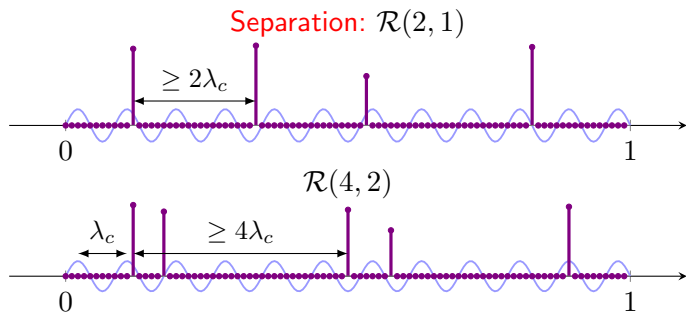
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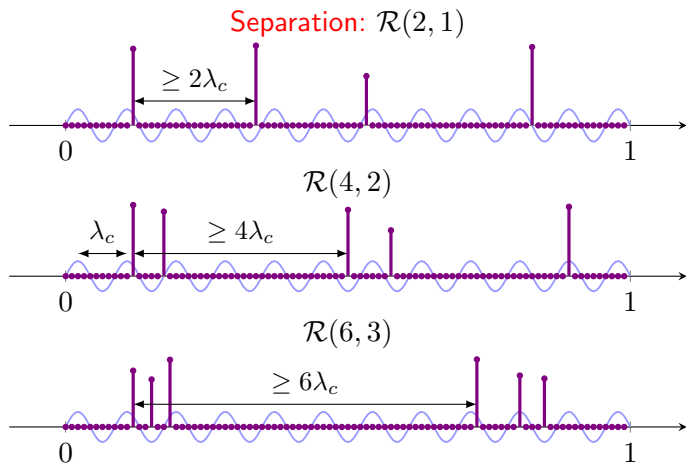
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\mathbf{x} has fewer than r spikes in every $\lambda_c d$ interval [$\lambda_c \triangleq 1/f_c$]

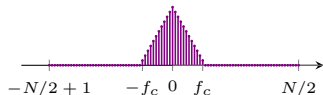


Main results

Recall:

$$\mathbf{s} = \mathbf{P}\mathbf{x} + \mathbf{z}$$

spectrum



Solve:

$$\text{minimize } \|\mathbf{s} - \mathbf{P}\hat{\mathbf{x}}\|_1 \quad \text{subject to } \hat{\mathbf{x}} \geq 0$$

Theorem: (V.Morgenshter, Càndes'14, [1])

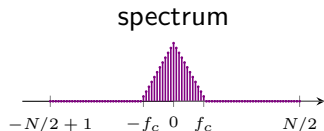
Take $\mathbf{P} = \mathbf{P}_{\text{tri}}$ or $\mathbf{P} = \mathbf{P}_{\text{flat}}$. Assume $\mathbf{x} \geq 0$, $\mathbf{x} \in \mathcal{R}(2r, r)$. Then,

$$\|\hat{\mathbf{x}} - \mathbf{x}\|_1 \leq c \cdot \|\mathbf{z}\|_1 \cdot \left(\frac{N}{2f_c}\right)^{2r}.$$

Main results

Recall:

$$\mathbf{s} = \mathbf{P}\mathbf{x} + \mathbf{z}$$



Solve:

$$\text{minimize } \|\mathbf{s} - \mathbf{P}\hat{\mathbf{x}}\|_1 \quad \text{subject to } \hat{\mathbf{x}} \geq 0$$

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$$\|\hat{\mathbf{x}} - \mathbf{x}\|_1 \leq c \cdot \|\mathbf{z}\|_1 \cdot \left(\frac{N}{2f_c}\right)^{2r}.$$

Converse: (V.Morgenshter, Càndes'14, [1])

For $\mathbf{P} = \mathbf{P}_{\text{tri}}$, no algorithm can do better than $c \cdot \|\mathbf{z}\|_1 \cdot \left(\frac{N}{2f_c}\right)^{2r-1}$.

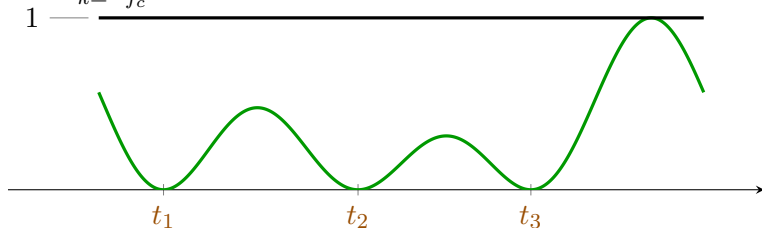
Key ideas

- **Duality theory:** to prove stability we need a low-frequency trigonometric polynomial that is “curvy”
 - [Dohono, et al.’92, Fuchs’05] construct trigonometric polynomial that is not “curvy”
 - [Candès and Fernandez-Granda’12] construct trigonometric polynomial that is “curvy”, but construction needs separation
 - **New construction:** multiply “curvy” trigonometric polynomials
 - “curvy”
 - construction needs no separation

Dual certificate (noisy case)

- \mathcal{T} is the support of \mathbf{x}
- Suppose, we can construct a **low-frequency trig. polynomial**:

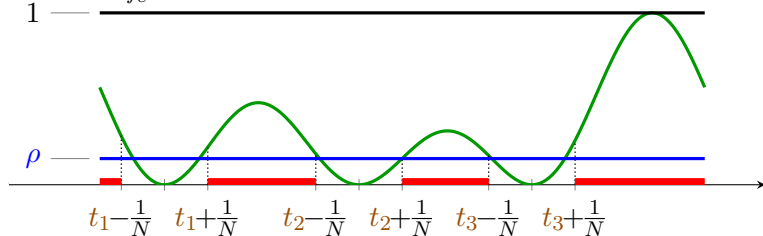
$$q(t) = \sum_{k=-f_c}^{f_c} \hat{q}_k e^{-i2\pi kt}, \quad 0 \leq q(t) \leq 1, \quad q(t_i) = 0 \text{ for all } t_i \in \mathcal{T}.$$



Dual certificate (noisy case)

- \mathcal{T} is the support of \mathbf{x}
- Suppose, we can construct a **low-frequency trig. polynomial**:

$$q(t) = \sum_{k=-f_c}^{f_c} \hat{q}_k e^{-i2\pi kt}, \quad 0 \leq q(t) \leq 1, \quad q(t_i) = 0 \text{ for all } t_i \in \mathcal{T}.$$



- Then, $\|\hat{\mathbf{x}} - \mathbf{x}\|_1 \leq 4\|\mathbf{z}\|_1/\rho$.

Proof of Lemma

■ Set:

$$\mathbf{h} = [h_0 \cdots h_{N-1}]^T = \hat{\mathbf{x}} - \mathbf{x}, \quad \mathcal{T} = \{l/N : h_l < 0\}$$

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$$\begin{aligned} |\langle \mathbf{q} - \rho/2, \mathbf{h} \rangle| &= |\langle \mathbf{P}(\mathbf{q} - \rho/2), \mathbf{h} \rangle| = |\langle \mathbf{q} - \rho/2, \mathbf{P}\mathbf{h} \rangle| \\ &\leq \|\mathbf{q} - \rho/2\|_{\infty} \|\mathbf{P}\mathbf{h}\|_1 \leq \|\mathbf{P}\mathbf{x} - \mathbf{s} + \mathbf{s} - \mathbf{P}\hat{\mathbf{x}}\|_1 \\ &\leq \|\mathbf{P}\mathbf{x} - \mathbf{s}\|_1 + \|\mathbf{s} - \mathbf{P}\hat{\mathbf{x}}\|_1 \\ &\leq 2\|\mathbf{P}\mathbf{x} - \mathbf{s}\|_1 \leq 2\|\mathbf{z}\|_1. \end{aligned}$$

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$$|\langle \mathbf{q} - \rho/2, \mathbf{h} \rangle| = \left| \sum_{l=0}^{N-1} (q_l - \rho/2) h_l \right| = \sum_{l=0}^{N-1} (q_l - \rho/2) h_l \geq \rho \|\mathbf{h}\|_1 / 2.$$

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- Combining: $\|\mathbf{h}\|_1 \leq 4\|\mathbf{z}\|_1 / \rho.$

Key ideas

- **Duality theory:** to prove stability we need a low-frequency trigonometric polynomial that is “curvy”
- [Dohono, et al.’92, Fuchs’05] construct trigonometric polynomial that is not “curvy”
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- **New construction:** multiply “curvy” trigonometric polynomials
 - “curvy”
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Donoho, et al.'92, Fuchs'05, [2, 3]: "Classical" $q(t)$

$$q(t) = \prod_{t_0 \in \mathcal{T}} \frac{1}{2} [\cos(2\pi(t + 1/2 - t_0)) + 1].$$

Euler's formula:

$$\cos(2\pi t) = \frac{e^{i2\pi t} + e^{-i2\pi t}}{2}$$

Sparsity implies $q(t)$ is low-frequency:

$$q(t) = \sum_{k=-f_c}^{f_c} \hat{q}_k e^{-i2\pi kt} \text{ if } |\mathcal{T}| \leq f_c$$

$$s \leq |\mathcal{T}| \leq f_c = \frac{1}{2} \times \text{number of measurements}$$

No square-root bottleneck!

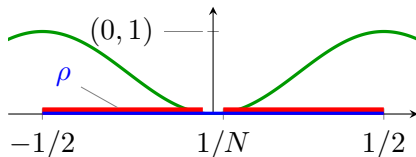
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No separation required

Low curvature!

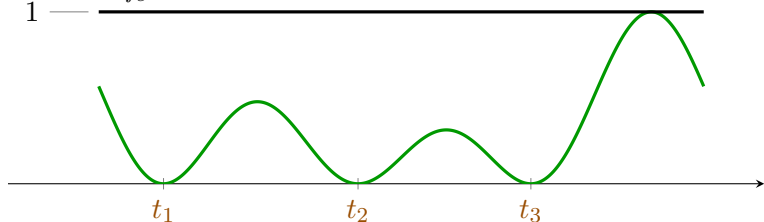
$$q(t - t_0) \approx (t - t_0)^2 \Rightarrow \|\mathbf{x} - \hat{\mathbf{x}}\|_1 \leq \|\mathbf{z}\|_1 \cdot N^2$$



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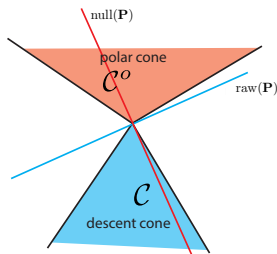


- Then, $\hat{\mathbf{x}} = \mathbf{x}$.

Connection to LASSO (\mathbf{x} can be negative here)

minimize $\|\hat{\mathbf{x}}\|_1$ subject to $\mathbf{s} = \mathbf{P}\hat{\mathbf{x}}$

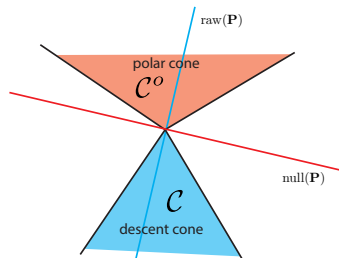
- $\hat{\mathbf{x}} = \mathbf{x}$ iff there exists $\mathbf{q} \perp \text{null}(\mathbf{P})$ and $\mathbf{q} \in \partial\|\mathbf{x}\|_1$
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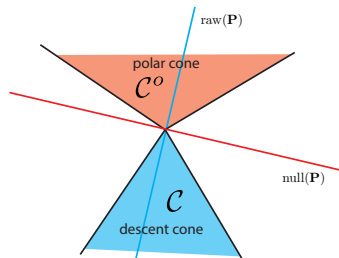
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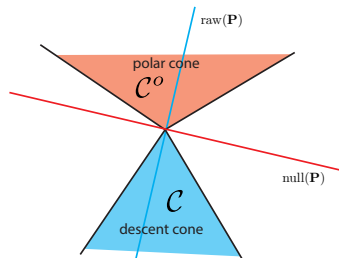
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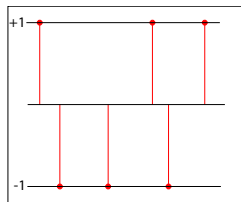
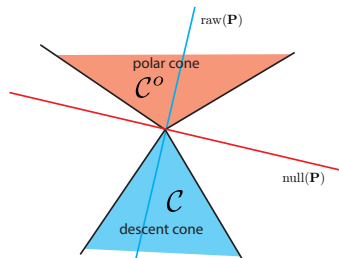
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$\text{sign}(x) \quad (x \neq 0)$

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Candès, Fernandez-Granda'12, [4]: “Curvy” $q(t)$

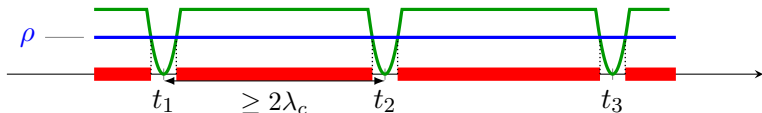
$$q(t) = \sum_{t_j \in \mathcal{T}} a_j K(t - t_j) + \text{corrections},$$

$K(t)$... low-frequency and “curvy”

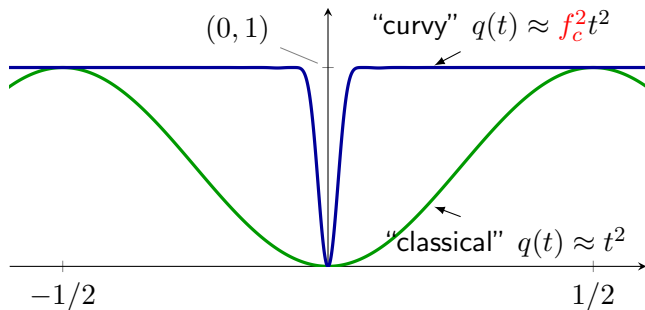
Separation between zeros required: $\mathcal{T} \in \mathcal{R}(2, 1)$

High curvature!

$$q(t - t_i) \approx f_c^2 (t - t_i)^2 \Rightarrow \|\mathbf{x} - \hat{\mathbf{x}}\|_1 \leq c \cdot \|\mathbf{z}\|_1 \cdot \left(\frac{N}{2f_c}\right)^2$$



Comparison of Trigonometric Polynomials



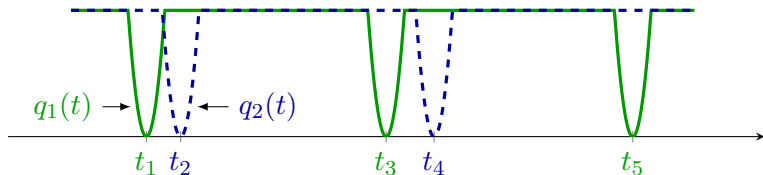
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New construction: curvature without separation

Partition support: $\mathcal{T} = \mathcal{T}_1 \cup \mathcal{T}_2$, $r = 2$

Regularity: $\mathcal{T} \in \mathcal{R}(2 \cdot 2, 2) \Rightarrow \mathcal{T}_i \in \mathcal{R}(4, 1)$

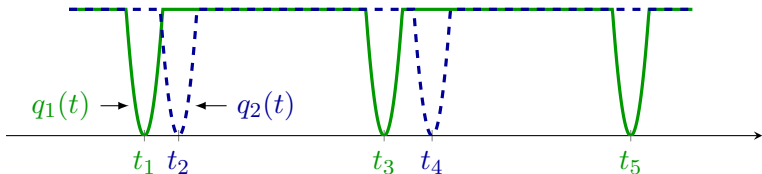


$$q(t; f_c) = q_1(t; f_c/2) \times q_2(t; f_c/2)$$

New construction: curvature without separation

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Summation vs. multiplication

Remember: $q(t)$ must be frequency-limited to f_c !

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This work:

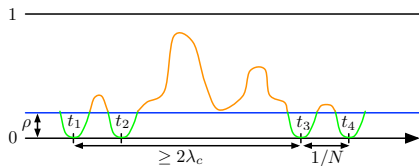
$$q(t) = \prod_{k=1}^r \sum_{t_{jk} \in \mathcal{T}_k} \underbrace{a_{jk} K(t - t_{jk})}_{\text{frequency } f_c/r}$$

Complex vs. positive signals

Why do we need $\mathbf{x} \geq \mathbf{0}$?

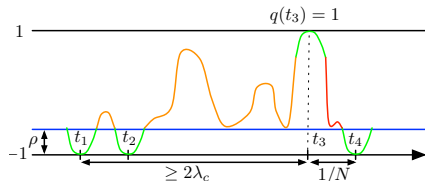
$$\mathbf{x} \geq \mathbf{0}$$

Interpolate **zero** on supp. of \mathbf{x}



$$\mathbf{x} \in \mathbb{C}^N$$

Interpolate $\text{sign}(\mathbf{x})$ on supp. of \mathbf{x}



Does not exist! (Bernstein Th.)

Continuous setup

f_c fixed, $N \rightarrow \infty \Rightarrow \text{SRF}_{\text{OLD}} \rightarrow \infty$



Is the problem hopeless?

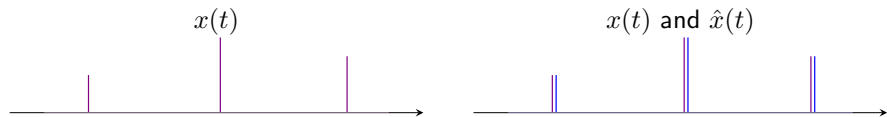
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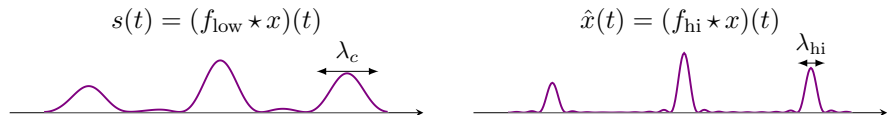
No: we need to be less ambitious!

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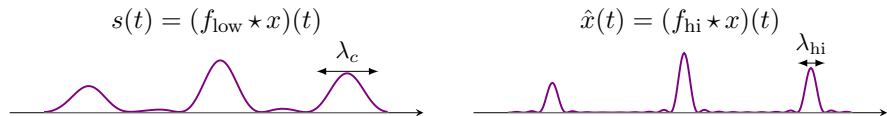


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$$\text{Error} = \|f_{\text{hi}} \star (x - \hat{x})\|_1$$



$$\text{SRF}_{\text{NEW}} = \lambda_c / \lambda_{\text{hi}}$$

Theorem: (V. Morgenshtern, 2019, [5])

Assume $x(t) \geq 0$, $x(t) \in \mathcal{R}(2r, r)$. Then,

$$\|f_{\text{hi}} \star (x - \hat{x})\|_1 \leq c \cdot \left(\frac{\lambda_c}{\lambda_{\text{hi}}}\right)^{2r} \cdot \|z(t)\|_1.$$

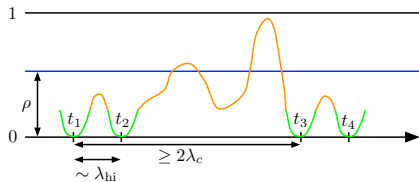
Need new tools

Theorem: (V. Morgenshtern, 2019, [5])

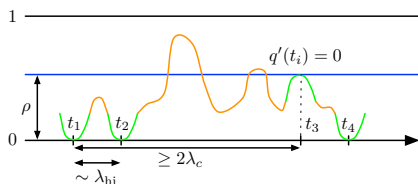
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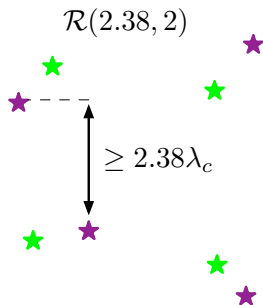
Can do: all zeros



Need: arbitrary pattern $\{0, +\rho\}$



2D Super-resolution



Theorem: (V. Morgenshtern and E. Candès, 2016, [1])

Take $\mathbf{P} = \mathbf{P}_{\text{tri},2\text{D}}$ or $\mathbf{P} = \mathbf{P}_{\text{flat},2\text{D}}$. Assume $\mathbf{x} \geq 0$, $\mathbf{x} \in \mathcal{R}(2.38r, r)$. Then,

$$\|\hat{\mathbf{x}} - \mathbf{x}\|_1 \leq c \cdot \left(\frac{N}{2f_c}\right)^{2r} \delta.$$

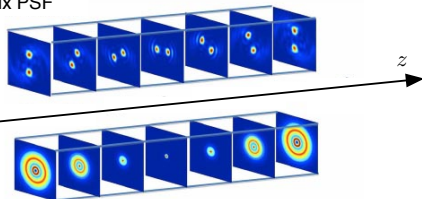
New: number of spikes is linear in the number of observations

Improving microscopes

Collaboration with Moerner Lab, C.A. Sing-Long, E. Candès

Reconstruction of 3D signals from 2D data

Double-helix PSF



Normal PSF

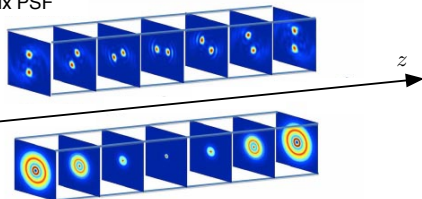
picture from [Pavani and Piston'08]



2D double-helix data

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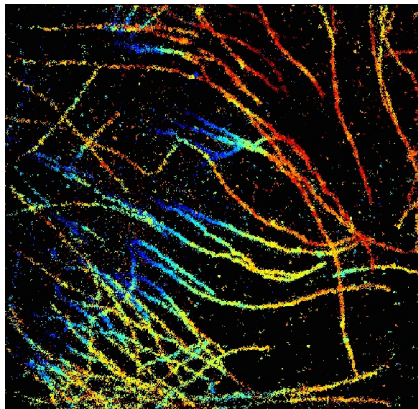


2D double-helix data

$$\begin{aligned} & \text{minimize} && \frac{1}{2} \|\mathbf{s} - \mathbf{P}\hat{\mathbf{x}}\|_2^2 + \lambda \sigma \|\text{diag}(\mathbf{w})\hat{\mathbf{x}}\|_1 \\ & \text{subject to} && \hat{\mathbf{x}} \geq 0 \end{aligned}$$

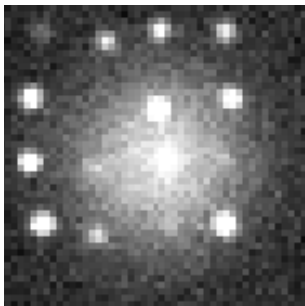
\mathbf{P} contains double-helix PSF slices

Preliminary result: 4 times faster than state-of-the-art



10000 CVX problems solved
TFOCS first order solver
millions of variables

Flexible framework: smooth background separation



minimize $\frac{1}{2} \|\mathbf{s} - \mathbf{P}(\hat{\mathbf{x}} + \mathbf{b})\|_2^2 + \lambda \sigma \|\hat{\mathbf{x}}\|_1$
subject to $\hat{\mathbf{x}} \geq 0$
 \mathbf{b} low freq. trig. polynomial (background)

Convex optimization is a near-optimal method
for super-resolution of positive sources

- Flexibility and good practical performance
- Non-asymptotic precise stability bounds
- Rayleigh-regularity is fundamental: separation between spikes is only one part of the picture

Backup slides

Connection to Bernstein theorem

Consider: $q(t) = \sum_{k=-f_c}^{f_c} \hat{q}_k e^{-i2\pi kt}$ with $\|q\|_\infty \leq 1$

Then: $\|q'\|_\infty \leq 2f_c$

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“Curvy” $q(t)$ has best possible curvature!

Since

$$q(t_i) = 0$$

$$q'(t_i) = 0$$

$$\|q\|_\infty \leq 1$$

We conclude:

$$\|q'\|_\infty \leq 2f_c \Rightarrow \|q''\|_\infty \leq (2f_c)^2$$

$$\Rightarrow q(t - t_i) \leq (2f_c)^2 (t - t_i)^2$$

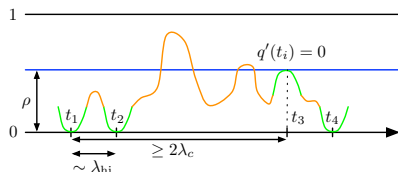
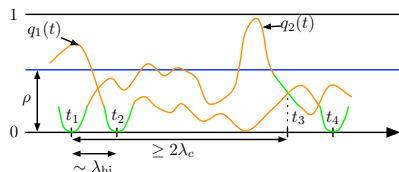
$$\Rightarrow q(t_i + 1/N) \leq \frac{(2f_c)^2}{N^2} = \frac{1}{\text{SRF}^2}$$

New tools

- 1 Control behavior on **separated set**
- 2 **Multiply**

$$q(t) = q_1(t) \times q_2(t)$$

$$0 = q'(t_3) = q_1'(t_3)q_2(t_3) + q_1(t_3)q_2'(t_3)$$

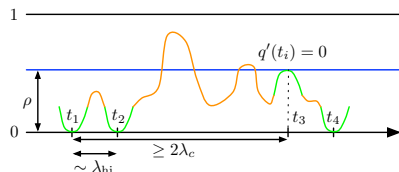
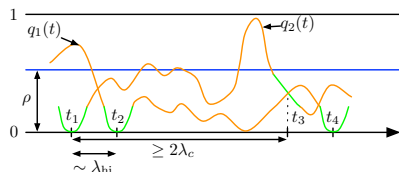


New tools

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- 3 **Sum**

$$q(t) = \sum_r \prod_{k=1}^r \sum_{t_{jk} \in \mathcal{T}_k} \underbrace{a_{jk} K(t - t_{jk})}_{\text{frequency } f_c/r}$$

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To be submitted.